

# Stochastic Model Predictive Control

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# Stochastic MPC

- Algorithmic considerations
- Implementation issues

# Outline

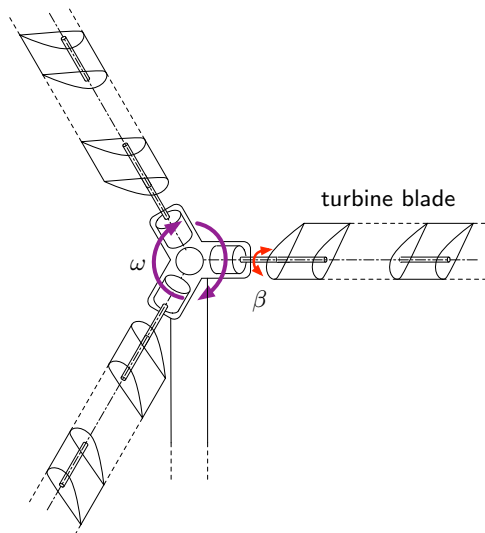
- ① Motivation: wind turbine blade fatigue control
- ② Stochastic MPC: basic formulations
  - ▶ Probabilistic constraints & recursive feasibility
  - ▶ Performance costs and stability analyses
- ③ Implementation
  - ▶ Affine model uncertainty: approximate and exact tubes
  - ▶ Additive model uncertainty: exact tubes

# Wind turbine blade pitch control

- System model
- High-cycle fatigue
- Constraints

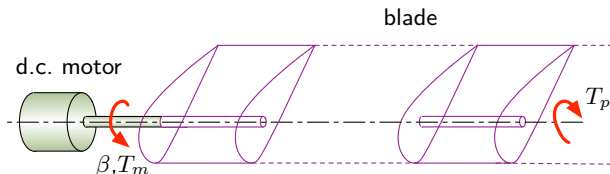
# Wind turbine blade pitch control

- Wind speed:  $v$   
assumed random
- Rotational speed:  $\omega$   
assumed constant
- Blade pitch angle:  $\beta$
- Power-optimal  
pitch angle:  $\beta^*(v, \omega)$



Control problem: track  $\beta^*$  subject to constraints on fatigue damage

# Blade dynamic model



blade pitch angle:  $\beta$

motor torque:  $T_m$

aerodynamic torque:  $T_p$

friction torque:  $c \frac{d\beta}{dt}$

Lumped parameter model

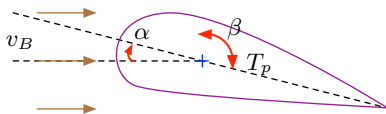
$$J \frac{d^2 \beta}{dt^2} + c \frac{d\beta}{dt} = T_m - T_p$$

# Aerodynamic torque

$$T_p = T_p(\alpha, v_B)$$

wind speed at blade:  $v_B = v_B(\omega, v)$

angle of attack:  $\alpha = \alpha(\beta, \omega, v)$



Hence the model

$$J \frac{d^2 \beta}{dt^2} + c \frac{d\beta}{dt} = T_m - T_p(\alpha, v_B)$$

contains:

**multiplicative** uncertainty due to  $v_B(\omega, v)$

**additive** uncertainty due to  $\alpha(\beta, \omega, v)$

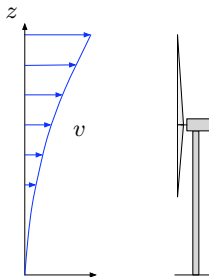
} statistically dependent

# Wind model

Wind speed variation with height  $z$ :

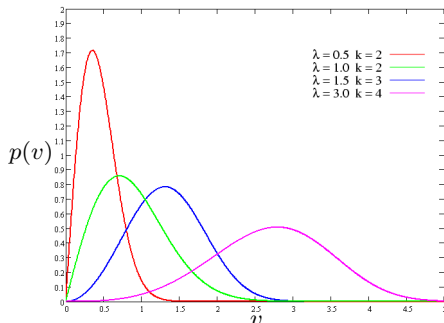
$$v = v_{\text{ref}} \frac{\log(z/z_0)}{\log(z_{\text{ref}}/z_0)}$$

(ground roughness factor:  $z_0 > 0$ )



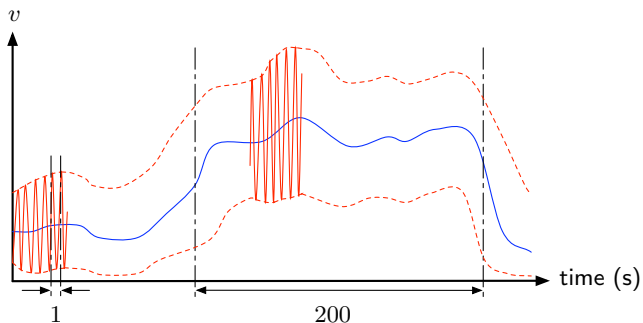
Wind speed pdf  $p(v)$

modelled using Weibull  
or Gaussian distributions





# Wind model



— :  $v$   
approx. constant mean and variance over interval  $\sim 200$  s

— :  $v_B$   
fluctuates due to blade rotation with period  $\sim 1$  s

# Linearized discrete-time system model

- Control **input**  $u$  determines motor torque:  $T_m = k_m u$
- System **output**: net torque  $y = T_m - T_p$
- Approximate linear model around a setpoint identified as:

$$\beta_{k+1} = a_{k,1}\beta_k + a_{k,0}\beta_{k-1} + b_{k,1}u_k + b_{k,0}u_{k-1} + \gamma_k$$

using:

NACA 632-215(V) blade data

1 second sampling interval

least squares estimation of  $(\bar{\theta}, \Sigma_\theta)$ ,  $\theta = [a_1 \ a_2 \ b_1 \ b_2 \ \gamma]^T$

# Linearized discrete-time system model

- Model:  $x(k+1) = A_k x(k) + B_k u(k) + w_k$ ,

$$A_k = \begin{bmatrix} 0 & a_{k,2} \\ 1 & a_{k,1} \end{bmatrix}, \quad B_k = \begin{bmatrix} b_{k,2} \\ b_{k,1} \end{bmatrix}, \quad w_k = \begin{bmatrix} 0 \\ \gamma_k \end{bmatrix}$$

with parameters

$$\begin{bmatrix} A_k & w_k \end{bmatrix} = \begin{bmatrix} \bar{A} & 0 \end{bmatrix} + \sum_{j=1}^3 \begin{bmatrix} A^{(j)} & w^{(j)} \end{bmatrix} q_j(k)$$

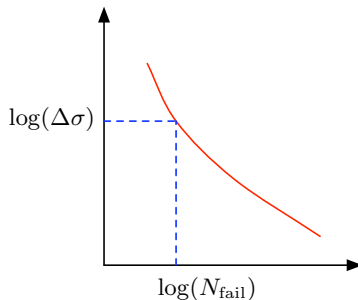
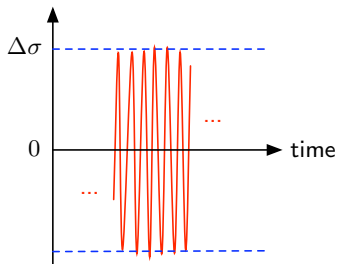
- Output:  $y(k) = c_k x(k) + d_k$ :

$$\begin{bmatrix} c_k & d_k \end{bmatrix} = \begin{bmatrix} \bar{c} & 0 \end{bmatrix} + \sum_{j=1}^2 \begin{bmatrix} c^{(j)} & d^{(j)} \end{bmatrix} q_j(k)$$

- Random variable  $q(k) \sim \mathcal{D}$ , approximately i.i.d.,  
identified empirically as truncated Gaussian

# High-cycle fatigue

- S-N curve gives number of cycles to failure ( $N_{\text{fail}}$ ) under cyclical stress loading:



- Extend to: **non-zero mean stress** (using Goodman's rule)  
**combined stress loadings** (using Miner's rule)

# Constraints

- Blade stress amplitude depends on net torque,  $y$ :

$$y = T_m - T_p$$

hence for a given life-span  $N_{\text{life}}$ , require:

$$N_{\text{fail}}/N_{\text{life}} \geq \text{rate of violation of bound: } y \leq \bar{y}$$

for  $\bar{y}$  and  $N_{\text{fail}}$  obtained from S-N curves

- Equivalent **probabilistic** constraints:

$$\Pr\{c_k x(k) + d_k \leq \bar{y}\} \geq p, \quad p = 1 - \frac{N_{\text{fail}}}{N_{\text{life}}}$$

- Motor torque saturation:  $|T_m| \leq T_{\text{max}}$  implies **hard** input constraints

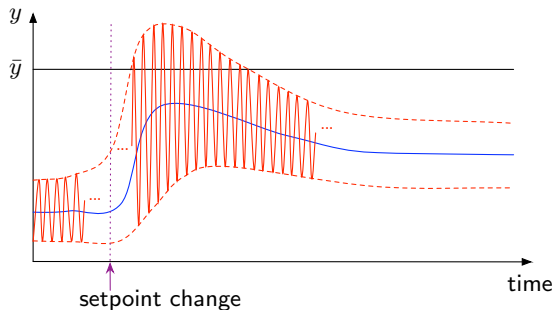
$$|u(k)| \leq \bar{u}$$

# Constraints

More accurate probabilistic constraints

... on the average rate of violation of bounds over interval  $N_s$

$$\frac{1}{N_s} \sum_{k=0}^{N_s-1} \Pr\{c_k x(k) + d_k \leq \bar{y}\} \geq p, \quad p = 1 - \frac{N_{\text{fail}}}{N_{\text{life}}}$$



... on ranges of stress amplitudes, e.g.  $\Pr\{c_k x(k) + d_k > \bar{y}_r\} < p_r = N_{\text{fail},r}/N_{\text{life}}$

$\vdots$

$$\Pr\{\bar{y}_2 \leq c_k x(k) + d_k > \bar{y}_1\} < p_1 = N_{\text{fail},1}/N_{\text{life}}$$

## Problem formulation

- System and constraints:

$$\begin{aligned}
 x^+ &= Ax + Bu + w, & A &= A(q), \quad B = B(q), \quad w = w(q) \\
 \Pr\{cx + d \leq \bar{y}\} &\geq p, & c &= c(q), \quad d = d(q) \\
 |u| &\leq \bar{u}
 \end{aligned}$$

Uncertainty:  $q \sim \mathfrak{D}$ , where distribution  $\mathfrak{D}$  is **finitely supported** in practice

- MPC law  $u(k) = \kappa_{\text{MPC}}(x(k))$  obtained by solving, at  $k = 0, 1, \dots$ :

$$\begin{aligned}
 \min_{\{u(k), u(k+1), \dots\}} & J(x(k), \{u(k), u(k+1), \dots\}) \\
 \text{subject to} & \quad \Pr\{c(k+i)x(k+i) + d(k+i) \leq \bar{y}\} \geq p \\
 & \quad |u(k+i)| \leq \bar{u}, \quad i = 0, 1, \dots
 \end{aligned}$$

How can we ensure recursive feasibility?

analyse closed loop behaviour?

invoke probabilistic constraints?

parameterize predictions?

## Probabilistic constraints & recursive feasibility



## Recursive feasibility

Consider the general dynamics:  $x^+ = f(x, u, q)$

and probabilistic constraint:  $\Pr\{F(x, u, q) \leq 1\} \geq p, \quad q \sim \mathfrak{D}$

Suppose  $u(0)$  and  $u(1)$  are such that

$$\Pr\{F(x(k), u(k), q(k)) \leq 1 \mid x(0)\} \geq p, \quad k = 0, 1$$

then:

(i) it is not necessarily true that

$$\Pr\{F(x(k), u(k), q(k)) \leq 1 \mid x(1)\} \geq p, \quad k = 1$$

e.g. take  $x \in \mathbb{R}$ ,  $F(x, u, q) = f(x, u, q)$ :

$$\Pr\{F(x(1), u(1), q(1)) \leq 1 \mid x(0)\} \geq p$$

but

$$\Pr\{F(x(1), u(1), q(1)) \leq 1 \mid x_{\max}(1)\} \not\geq p$$

for some realizations  $q(0)$ , there may not exist any  $u(1)$  satisfying this constraint

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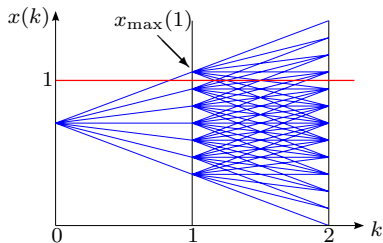
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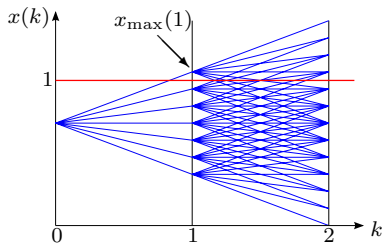
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(ii) for some realizations  $q(0)$ , there may not exist **any**  $u(1)$  satisfying this constraint

# Recursive feasibility

Define  $\mathcal{T}_0 = \{x : \exists u, F(x, u, q) \leq 1 \text{ w.p. } p\}$ ,

then  $u(1)$  exists s.t.  $F(x, u, q) \leq 1 \text{ w.p. } p$  iff  $x(1) \in \mathcal{T}_0 \text{ w.p. } 1 \leftarrow$  **hard constraint**

$$\exists \{u(0), u(1), \dots\} \text{ s.t. } \left\{ \begin{array}{l} F(x(k), u(k), q) \leq 1 \text{ w.p. } p \\ x(k) \in \mathcal{T}_0 \text{ w.p. } 1 \end{array} \right\} k = 0, 1, \dots \text{ iff } x(0) \in \mathcal{R}_\infty$$

$\mathcal{R}_\infty =$  infinite-time reachability set

[Bertsekas 1972]

$$\hat{\mathcal{T}}_0 = \{(x, u) : F(x, u, q) \leq 1 \text{ w.p. } p\},$$

$$\mathcal{R}_0 = \mathcal{T}_0 = \text{Proj}_x(\hat{\mathcal{T}}_0)$$

$$\hat{\mathcal{R}}_k = \{(x, u) : f(x, u, q) \in \mathcal{R}_{k-1} \text{ w.p. } 1\} \cap \hat{\mathcal{T}}_0, \quad \mathcal{R}_k = \text{Proj}_x(\hat{\mathcal{R}}_k), \quad k = 1, 2, \dots$$

$\mathcal{R}_\infty$  exists if  $\mathcal{T}_0$  and  $\text{supp}(\mathcal{D})$  are compact, and  $x^+ = f(x, u, q)$  is stabilizable, with minimal robust control invariant set contained in  $\text{int}(\mathcal{T}_0)$ .

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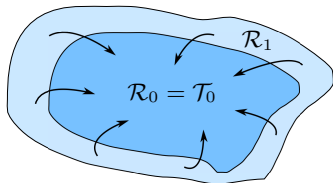
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# Terminal set

Let  $\mathcal{O}_\infty =$  maximal admissible set for dynamics:  $x^+ = f(x, \kappa_{\mathbb{T}}(x), q)$   
and constraint:  $F(x, \kappa_{\mathbb{T}}(x), q) \leq 1$  w.p.  $p$

where  $\mathcal{O}_k = \bigcap_{j=0}^k \mathcal{S}_j$ ,  $k = 0, 1, \dots$

[cf. Kolmanovsky & Gilbert 1998]

$$\mathcal{S}_0 = \{x : F(x, \kappa_{\mathbb{T}}(x), q) \leq 1 \text{ w.p. } p\}$$

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If  $\mathcal{O}_k \subseteq \mathcal{S}_{k+1}$ , then  $\mathcal{O}_k = \mathcal{O}_\infty$

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$$\mathcal{S}_k = \{x(0) : x(k) \in \mathcal{S}_0 \text{ w.p. } 1\}$$

If  $\mathcal{O}_k \subseteq \mathcal{S}_{k+1}$ , then  $\mathcal{O}_k \cap \mathcal{S}_{k+1} = \mathcal{O}_{k+1} = \mathcal{O}_k$ , and

$$(i) \quad x(0) \in \mathcal{O}_k = \mathcal{O}_{k+1} \Rightarrow \left\{ \begin{array}{c} x(0) \in \mathcal{S}_1 \\ \vdots \\ x(0) \in \mathcal{S}_{k+1} \end{array} \right\} \Rightarrow \left\{ \begin{array}{c} x(1) \in \mathcal{S}_0 \text{ w.p. } 1 \\ \vdots \\ x(1) \in \mathcal{S}_k \text{ w.p. } 1 \end{array} \right\} \Rightarrow x(1) \in \mathcal{O}_k \text{ w.p. } 1$$

i.e.  $\mathcal{O}_k$  is positively invariant with probability 1

$$(ii) \quad x(0) \in \mathcal{O}_k \Rightarrow x(i) \in \mathcal{O}_k \quad \forall i \geq 0 \Rightarrow x(i) \in \mathcal{S}_0 \text{ w.p. } 1 \quad \forall i \geq 0 \\ \Rightarrow x(0) \in \mathcal{S}_i \quad \forall i \geq 0$$

hence  $\mathcal{O}_k \subseteq \bigcap_{j=0}^{\infty} \mathcal{S}_j = \mathcal{O}_\infty$

but  $\mathcal{O}_\infty \subseteq \mathcal{O}_k$  by definition, so  $\mathcal{O}_k = \mathcal{O}_\infty$

# Recursive feasibility

Prototype stochastic MPC optimization:

$$\begin{array}{ll}
 \min_{\{u(k), u(k+1), \dots\}} & J(x(k), \{u(k), u(k+1), \dots\}) \\
 \text{subject to} & \left. \begin{array}{l} F(x(k+i), u(k+i), q(k+i)) \leq 1 \text{ w.p. } p \\ x(k+i) \in \mathcal{T}_0 \text{ w.p. } 1 \\ x(k+N) \in \mathbb{T} \text{ w.p. } 1 \end{array} \right\} i = 0, \dots, N-1
 \end{array}$$

- ★ Choose  $\mathbb{T}$  as the maximal admissible set  $\mathcal{O}_\infty$  (or an invariant inner approximation)
- ★ Recursive feasibility guarantee:
  - if  $\{u(0), \dots, u(N-1)\}$  is feasible for given  $x(0)$
  - then  $\{u(1), \dots, u(N-1), \kappa_{\mathbb{T}}(x(N))\}$  is feasible for  $x(1) = f(x(0), u(0), q)$  w.p. 1
- ★ Definition of cost  $J$  determines stability properties. Consider expected and nominal quadratic costs.

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 & \forall q(k+j) \in \text{supp}(\mathfrak{D}), j = 0, \dots, i-1 \\
 & x(k+N) \in \mathbb{T} \\
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- ★ Definition of cost  $J$  determines stability properties.  
Consider expected and nominal quadratic costs.

## Recursive feasibility

- Re-parameterize predicted control law:  $u = \kappa_{\mathbb{T}}(x) + c$   
then predicted control inputs,  $\{u(0), u(1), \dots\}$ , at time 0 are defined by

$$u(k) = \kappa_{\mathbb{T}}(x(k)) + c(k) \quad \text{with } c(k) = 0, \quad \forall k \geq N$$

$$\mathbf{c}(0) = \{c(0), \dots, c(N-1)\}$$

- Define feasible set  $\mathcal{F} = \{(x, \mathbf{c}) : \text{constraints are satisfied}\}$ , i.e.

$$\mathcal{F} = \left\{ \begin{array}{l} F(x(k+i), u(k+i), q(k+i)) \leq 1 \text{ w.p. } p \\ (x(k), \mathbf{c}(k)) : x(k+i) \in \mathcal{T}_0 \text{ w.p. } 1 \\ x(k+N) \in \mathbb{T} \text{ w.p. } 1 \end{array} \right\} \quad i = 0, \dots, N-1$$

and  $\mathcal{F}_x = \text{Proj}_x(\mathcal{F})$

Then

$$x(0) \in \mathcal{F}_x \quad \implies \quad (x(1), \tilde{\mathbf{c}}(1)) \in \mathcal{F} \text{ w.p. } 1$$

where  $x(1) = f(x(0), \kappa_{\mathbb{T}}(x(0)) + c(0), q)$

$$\tilde{\mathbf{c}}(1) = \{c(1), \dots, c(N-1), 0\}$$

## Performance cost, stability and convergence



# Local Stability

Lyapunov (robust) stability  $\rightsquigarrow$  local property of optimal unconstrained MPC law  $\kappa_{\text{MPC}}^{\text{uc}}$

Define

$\mathcal{M}_{\infty}^{\text{uc}}$  = minimal robust invariant set under  $\kappa_{\text{MPC}}^{\text{uc}}$

$$\mathcal{F}_x^{\text{uc}} = \left\{ x(0) : \begin{array}{l} (x(k), u(k)) \text{ satisfies constraints } \forall k \geq 0 \\ \text{under } u(k) = \kappa_{\text{MPC}}^{\text{uc}}(x(k)) \end{array} \right\}$$

★ If  $\kappa_{\text{MPC}}^{\text{uc}}$  is asymptotically stabilizing and  $\mathcal{M}_{\infty}^{\text{uc}} \subset \text{int}(\mathcal{F}_x^{\text{uc}})$ , then

$\mathcal{M}_{\infty}^{\text{uc}}$  is asymptotically stable

with region of attraction containing  $\mathcal{F}_x^{\text{uc}}$

★ If  $\kappa_{\mathbb{T}} =$  optimal unconstrained control law, then

$$\kappa_{\text{MPC}}^{\text{uc}} = \kappa_{\mathbb{T}} \text{ and } \mathcal{F}_x^{\text{uc}} \supseteq \mathbb{T}$$

and hence  $\mathcal{F}_x^{\text{uc}} \supseteq \mathbb{T} \supset \mathcal{M}_{\infty}^{\text{uc}}$

# Expectation MPC cost

Quadratic expected value cost:

$$J(x(k), \{u(k), u(k+1), \dots\}) = \sum_{i=0}^{\infty} \mathbb{E}(\|x(k+i)\|_Q^2 + \|u(k+i)\|_R^2)$$

optimal unconstrained control law:  $u(k) = K_{LQ}x(k)$

▷ Finite cost → minimize numerically online

▷ Quadratic in dof  $\mathbf{c}(k)$ :

$$V(x(k), \mathbf{c}(k)) = \mathbf{c}(k)^T P_{cc} \mathbf{c}(k) + 2x(k)^T P_{xc} \mathbf{c}(k) + p_k(x(k))$$

$P_{cc}, P_{xc}$  computed offline,  $P_{xc} = 0$  if  $\kappa_T = K_{LQ}$

▷ Optimal value  $V^*(x) = \min_{\mathbf{c} \in \mathcal{F}_c(x)} V(x, \mathbf{c})$  is lower-bounded in  $x$

▷ MPC law:  $\kappa_{\text{MPC}}(x) = \kappa_T(x) + \mathbf{c}^*$ , where  $\mathbf{c}^* = \arg \min_{\mathbf{c} \in \mathcal{F}_c(x)} V(x, \mathbf{c})$

# Expectation MPC cost

Equivalent MPC objective:

$$V(x(k), \mathbf{c}(k)) = \sum_{i=0}^{\infty} \left[ \mathbb{E}(\|x(k+i)\|_Q^2 + \|u(k+i)\|_R^2) - l_{ss} \right]$$

where  $l_{ss} = \lim_{k \rightarrow \infty} \mathbb{E}(\|x(k)\|_Q^2 + \|u(k)\|_R^2)$  under  $u(k) = K_{LQ}x(k)$

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$$V(x(k), \mathbf{c}(k)) = \mathbf{c}(k)^T P_{cc} \mathbf{c}(k) + 2x(k)^T P_{xc} \mathbf{c}(k) + p_k(x(k))$$

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# Expectation MPC cost: convergence analysis

Let

$$\begin{aligned} \mathbf{c}^*(0) &= \{c^*(0), \dots, c^*(N-1)\} = \arg \min_{\mathbf{c} \in \mathcal{F}_{\mathbf{c}}(x(0))} V(x(0), \mathbf{c}) \\ \tilde{\mathbf{c}}(1) &= \{c^*(1), \dots, c^*(N-1), 0\} \end{aligned}$$

Then, by definition, for  $x(0) \in \mathcal{F}_x$ :

$$V(x(1), \tilde{\mathbf{c}}(1)) = V^*(x(0)) - (\|x(0)\|_Q^2 + \|u(0)\|_R^2 - l_{ss})$$

but  $\tilde{\mathbf{c}}(1) \in \mathcal{F}_{\mathbf{c}}(x(1))$  w.p. 1, so the optimal value function satisfies

$$\mathbb{E}(V^*(x(1))) \leq V^*(x(0)) - (\|x(0)\|_Q^2 + \|u(0)\|_R^2 - l_{ss})$$

★ summing over  $r$  time-steps:

$$\frac{1}{r} \sum_{k=0}^{r-1} \mathbb{E}(\|x(k)\|_Q^2 + \|u(k)\|_R^2) \leq l_{ss} + \frac{1}{r} [V^*(x(0)) - \mathbb{E}(V^*(x(r)))]$$

★ hence

$$\lim_{r \rightarrow \infty} \frac{1}{r} \sum_{k=0}^{r-1} \mathbb{E}(\|x(k)\|_Q^2 + \|u(k)\|_R^2) \leq l_{ss}$$

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# Expectation MPC cost: convergence analysis

- If  $l_{ss} = 0$  (no additive disturbance), then

$$\mathbb{E}(V^*(x(1))) \leq V^*(x(0)) - (\|x(0)\|_Q^2 + \|u(0)\|_R^2)$$

and  $(x(k), u(k)) \rightarrow (0, 0)$  as  $k \rightarrow \infty$

- If  $\kappa_T = K_{LQ}$ , then  $l_{ss}$  is minimal:

$$l_{ss} = \min_{\{u(0), u(1), \dots\}} \lim_{r \rightarrow \infty} \frac{1}{r} \sum_{k=0}^{r-1} \mathbb{E}(\|x(k)\|_Q^2 + \|u(k)\|_R^2)$$

Hence

(i) the bound  $\lim_{r \rightarrow \infty} \frac{1}{r} \sum_{k=0}^{r-1} \mathbb{E}(\|x(k)\|_Q^2 + \|u(k)\|_R^2) \leq l_{ss}$  holds with equality

(ii)  $\lim_{k \rightarrow \infty} \kappa_{\text{MPC}}(x(k)) = K_{LQ}x(k)$

# Expectation MPC cost: convergence analysis

- If  $l_{\text{ss}} = 0$  (no additive disturbance), then

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- (ii)  $\lim_{k \rightarrow \infty} \kappa_{\text{MPC}}(x(k)) = K_{\text{LQ}}x(k)$



# Nominal MPC cost

Linear dynamics:  $x^+ = Ax + Bu + w$ ,  $(A, B, w) = (A(q), B(q), w(q))$ ,  $q \sim \mathcal{D}$   
 $\mathbb{E}(A, B, w) = (A^0, B^0, 0)$

State decomposition:  $x = z + e$

Control decomposition:  $u = v + K_{\top}e$ ,  $v = K_v z + c$

Nominal dynamics:  $z^+ = A^0 z + B^0 v$

▷ Nominal cost:

$$V_0(x(0), c(0), z(0)) = \sum_{k=0}^{\infty} (\|z(k)\|_Q^2 + \|v(k)\|_R^2)$$

▷ Optimal value:  $V_0^*(x) = \min_{(c, z) \in \mathcal{F}_{c, z}(x)} V_0(x, c, z)$  is lower-bounded in  $x$

▷ MPC law:  $\kappa_{\text{MPC}}(x) = K_{\top}e + v^*$ , where  $(c^*, z^*) = \arg \min_{(c, z) \in \mathcal{F}_{c, z}(x)} V_0(x, c, z)$   
 $v^* = K_v z^* + c^*$

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 $\mathbb{E}(A, B, w) = (A^0, B^0, 0)$

State decomposition:  $x = z + e$

Control decomposition:  $u = v + K_{\mathbb{T}}e$ ,  $v = K_v z + c$

Nominal dynamics:  $z^+ = A^0 z + B^0 v$

▷ Nominal cost:

$$V_0(x(0), \mathbf{c}(0), z(0)) = \sum_{k=0}^{\infty} (\|z(k)\|_Q^2 + \|v(k)\|_R^2)$$

▷ Optimal value:  $V_0^*(x) = \min_{(\mathbf{c}, z) \in \mathcal{F}_{\mathbf{c}, z}(x)} V_0(x, \mathbf{c}, z)$  is lower-bounded in  $x$

▷ MPC law:  $\kappa_{\text{MPC}}(x) = K_{\mathbb{T}}e + v^*$ , where  $(\mathbf{c}^*, z^*) = \arg \min_{(\mathbf{c}, z) \in \mathcal{F}_{\mathbf{c}, z}(x)} V_0(x, \mathbf{c}, z)$   
 $v^* = K_v z^* + c^*$

# Nominal MPC cost: convergence analysis

Recursive feasibility:

$$x(0) \in \mathcal{F}_x \implies (x(1), \tilde{\mathbf{c}}(1), \tilde{z}(1)) \in \mathcal{F} \text{ w.p. } 1$$

$$\text{where } x(1) = Ax(0) + B\kappa_{\text{MPC}}(x(0)) + w$$

$$\tilde{\mathbf{c}}(1) = \{c^*(1), \dots, c^*(N-1), 0\}$$

$$\tilde{z}(1) = A^0 z^*(0) + B^0 v^*(0)$$

Hence  $V_0(x(1), \tilde{\mathbf{c}}(1), \tilde{z}(1)) \leq V_0^*(x(0)) - (\|z^*(0)\|_Q^2 + \|v^*(0)\|_R^2)$ , and

$$V_0^*(x(1)) \leq V_0^*(x(0)) - (\|z^*(0)\|_Q^2 + \|v^*(0)\|_R^2)$$

★ sum over  $r$  time steps:

$$\sum_{k=0}^{r-1} (\|z^*(k)\|_Q^2 + \|v^*(k)\|_R^2) \leq V_0^*(x(0))$$

★ asymptotically:

$$\lim_{k \rightarrow \infty} (\|z^*(k)\|_Q^2 + \|v^*(k)\|_R^2) = 0 \implies \begin{cases} z^*(k) \rightarrow 0 \\ v^*(k) \rightarrow 0 \\ \kappa_{\text{MPC}}(x(k)) \rightarrow K_{\text{T}}x(k) \end{cases} \text{ as } k \rightarrow \infty$$

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Hence  $V_0(x(1), \tilde{\mathbf{c}}(1), \tilde{z}(1)) \leq V_0^*(x(0)) - (\|z^*(0)\|_Q^2 + \|v^*(0)\|_R^2)$ , and

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# Nominal MPC cost: convergence analysis

- If  $K_{\mathbb{T}}$  is mean-square stabilizing, i.e.  $\exists P_{\mathbb{T}} \succ 0$ :

$$P_{\mathbb{T}} - \mathbb{E}[(A + BK_{\mathbb{T}})^T P_{\mathbb{T}}(A + BK_{\mathbb{T}})] = I$$

then the closed loop dynamics under  $u = \kappa_{\text{MPC}}(x)$ :

$$x^+ = (A + BK_{\mathbb{T}})x + w', \quad w' = B(v^* - K_{\mathbb{T}}z^*) + w$$

has finite  $l^2$ -gain ( $w' \rightarrow x$ ),  $\gamma$ , i.e.  $\exists \gamma > 0$ :

$$\sum_{k=0}^{r-1} \mathbb{E}(\|x(k)\|^2) \leq \|x(0)\|_{P_{\mathbb{T}}}^2 + \gamma \sum_{k=0}^{r-1} \mathbb{E}(\|Bv^*(k) - K_{\mathbb{T}}z^*(k) + w(k)\|^2)$$

- Hence

$$\frac{1}{r} \sum_{k=0}^{r-1} \mathbb{E}(\|x(k)\|^2) \leq \gamma \mathbb{E}(\|w(0)\|^2) + \frac{1}{r} \|x(0)\|_{P_{\mathbb{T}}}^2 + \frac{\gamma'}{r} V_0^*(x(0))$$

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# Expectation vs Nominal MPC cost

Both expectation and nominal cost ensure: robust asymptotic stability of  $\mathcal{M}_\infty^{\text{uc}}$   
and finite  $l^2$ -gain ( $w \rightarrow x$ )

However

- ▷ nominal can be less representative of predicted performance than expectation cost

↑  
since  $e^*(0) = x(0) - z^*(0)$  is not necessarily zero

- ▷ nominal allows different (linear) feedback gains in nominal and disturbed dynamics

↑  
extra flexibility enables e.g.  $K_{\mathbb{T}}$  chosen for large  $\mathbb{T} = \mathcal{O}_\infty$   
 $K_v$  chosen for good  $l^2$  performance:  $K_v = K_{\text{LQ}}$

# Problem formulation: summary

Prototype stochastic MPC optimization:

$$\begin{array}{ll}
 \min_{\mathbf{c}} & \min_{\{u(k), u(k+1), \dots\}} J(x(k), \{u(k), u(k+1), \dots\}) \\
 & \text{subject to } \left. \begin{array}{l} F(x(k+i), u(k+i), q(k+i)) \leq 1 \text{ w.p. } p \\ x(k+i) \in \mathcal{T}_0 \text{ w.p. } 1 \\ x(k+N) \in \mathbb{T} \text{ w.p. } 1 \end{array} \right\} i = 0, \dots, N-1
 \end{array}$$

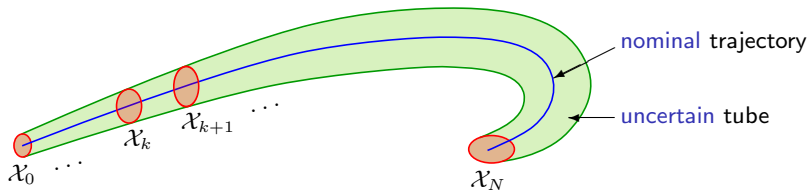
- ★ expectation or nominal cost
- ★ constraint at prediction time  $k+i$  invoked with:
  - worst-case uncertainty at prediction times  $k, k+1, \dots, k+i-1$
  - stochastic uncertainty at prediction time  $k+i$



# Propagating uncertain predictions

## Tube MPC

Split predicted trajectories (input/state) into **nominal** + **uncertain** components



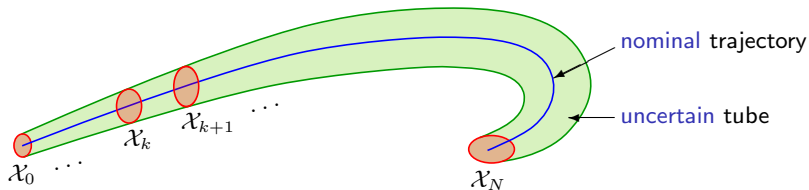
Applications to:

- ★ Linear model plus additive uncertainty or uncertain state estimation  
Langson ('04), Mayne ('05)
- ★ Nonlinear model plus uncertainty  
Blanchini ('90), Lee ('02), Raković ('06)
- ★ Linear and nonlinear stochastic systems with additive or multiplicative uncertainty  
Cannon ('09), Cannon ('10)

# Propagating uncertain predictions

## Tube MPC

Split predicted trajectories (input/state) into **nominal** + **uncertain** components



Computational advantages:

- ★ Offline computation assuming a fixed disturbance feedback law

- ▷ linear systems

Gossner ('97), Langson ('04), Mayne ('05)

- ★ Outer approximation using one-step-ahead predictions

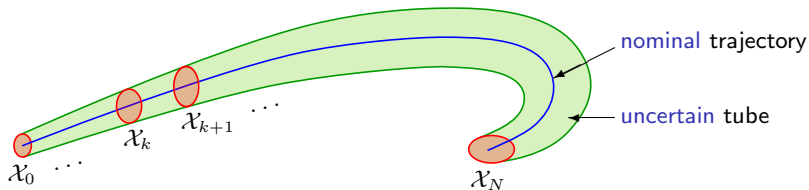
- ▷ nonlinear or stochastic systems

Blanchini ('90), Lee ('02), Cannon ('09)

# Propagating uncertain predictions

## Tube MPC

Split predicted trajectories (input/state) into **nominal** + **uncertain** components



Different approaches to computing probability distributions

Model uncertainty	Exact tubes	Approximate tubes
additive	numerical integration	numerical integration
additive and multiplicative	sampling/M-C simulation	parametric bounds

## Affine uncertainty: Approximate tubes

# Affine uncertainty: Parameter bounds

- Affine in the disturbance dynamics:

$$x_{k+1} = A(q_k)x_k + B(q_k)u_k + d(q_k)$$

$$(A, b, d) = (A^0, B^0, 0) + \sum_{i=1}^m (A^{(i)}, B^{(i)}, d^{(i)})q_{j,k}$$

$$q_k \sim \mathcal{D} \text{ i.i.d.}, \text{ supp}(\mathcal{D}) = \mathcal{Q}$$

- Confidence region: if  $q_k \in \hat{\mathcal{Q}}(p)$  w.p.  $p$ , where  $\hat{\mathcal{Q}}(p) = \text{Co}\{q^{(i)}, i = 1, \dots, l\}$ , then

$$x_{k+1} \in \text{Co}\{A(q^{(i)})x_k + B(q^{(i)})u_k + d(q^{(i)}), i = 1, \dots, l\} \text{ w.p. } p$$

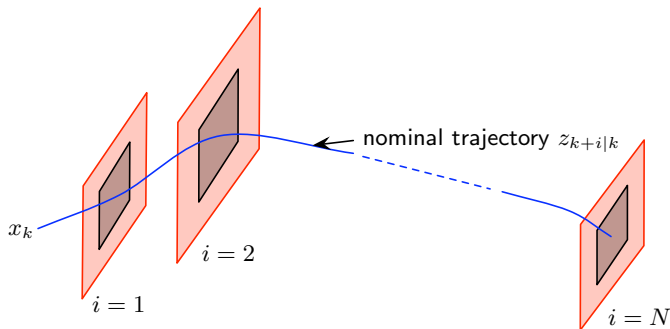
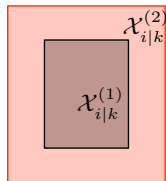
$$= \text{Co}\{\hat{A}^{(i)}x_k + \hat{B}^{(i)}u_k + \hat{d}^{(i)}, i = 1, \dots, l\}$$

# Affine uncertainty: Probabilistic tubes

Tube cross-section at prediction time-step  $i$ :

$$\{\mathcal{X}_{i|k}^{(1)}, \dots, \mathcal{X}_{i|k}^{(r)}\}, \text{ e.g. } \mathcal{X}_{i|k}^{(j)} = z_{k+i|k} + \{e : |e| \leq \bar{e}_{i|k}^{(j)}\}$$

$$\text{with } \mathcal{X}_{i|k}^{(1)} \subseteq \dots \subseteq \mathcal{X}_{i|k}^{(r)}$$

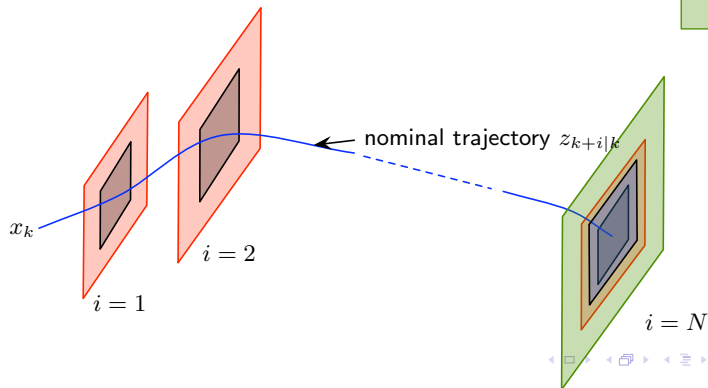
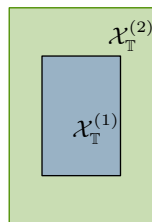


# Affine uncertainty: Probabilistic tubes

Terminal sets:

$$\{\mathcal{X}_{\mathbb{T}}^{(1)}, \dots, \mathcal{X}_{\mathbb{T}}^{(r)}\}, \quad \mathcal{X}_{\mathbb{T}}^{(j)} = \{x : |x| \leq \bar{x}_{\mathbb{T}}^{(j)}\}$$

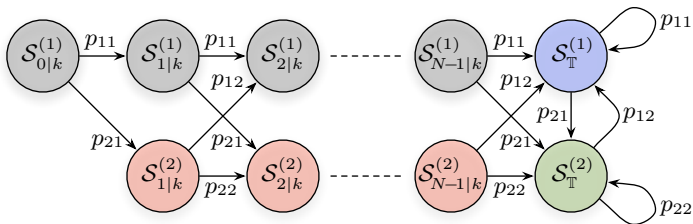
$$\text{with } \mathcal{X}_{\mathbb{T}}^{(1)} \subseteq \dots \subseteq \mathcal{X}_{\mathbb{T}}^{(r)}$$



## Affine uncertainty: Probabilistic tubes

$$\text{Let } \mathcal{S}_\star^{(j)} = \begin{cases} \mathcal{X}_\star^{(1)} & j = 1 \\ \mathcal{X}_\star^{(j)} - \mathcal{X}_\star^{(j-1)} & j = 2, \dots, r \end{cases}$$

(i). Define **transition probabilities**  $p_{jm}$ ,  $j, m = 1, \dots, r$ :



then  $p_i^{(j)} = \Pr(e_{k+i|k} \in \mathcal{S}_{i|k}^{(j)})$  is given by

$$\begin{bmatrix} p_i^{(1)} \\ p_i^{(2)} \\ \vdots \\ p_i^{(r)} \end{bmatrix} = \Pi^i \begin{bmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix}, \quad \Pi = \begin{bmatrix} p_{11} & \cdots & p_{1r} \\ \vdots & \ddots & \vdots \\ p_{r1} & \cdots & p_{rr} \end{bmatrix}$$



## Affine uncertainty: Probabilistic tubes

$$\text{Let } \mathcal{S}_\star^{(j)} = \begin{cases} \mathcal{X}_\star^{(1)} & j = 1 \\ \mathcal{X}_\star^{(j)} - \mathcal{X}_\star^{(j-1)} & j = 2, \dots, r \end{cases}$$

(ii). Define  $p_j$  as the **probability of satisfying soft constraints** in  $\mathcal{S}_\star^{(j)}$ :

$$\Pr(f^T x_{k+i+1|k} + g^T u_{k+i+1|k} > h \mid e_{k+i|k} \in \mathcal{S}_{i|k}^{(j)}) < p_j \quad i \leq N - 1$$

$$\Pr(f^T x_{k+i+1|k} + g^T u_{k+i+1|k} > h \mid x_{k+i|k} \in \mathcal{S}_\mathbb{T}^{(j)}) < p_j \quad i \geq N$$

(i) and (ii) imply

$$\Pr(f^T x_{k+i+1|k} + g^T u_{k+i+1|k} > h) < [p_1 \quad p_2 \quad \cdots \quad p_r] \Pi^i e_1$$

for all  $i$

## Affine uncertainty: Probabilistic tubes

The constraint

$$\Pr(f^T x_{k+i|k} + g^T u_{k+i|k} > h) < p$$

is satisfied for all  $i$  if  $\Pi$  and  $p_1, \dots, p_r$  satisfy

$$[p_1 \quad p_2 \quad \cdots \quad p_r] \Pi^i e_1 < p, \quad \forall i$$

- ★ Fix  $\Pi$  and  $p_1, \dots, p_r$  offline and optimize  $\{\mathcal{X}_{i|k}^{(j)}\}$  online subject to constraints on:
  - (i). transition probabilities
  - (ii). probabilities of satisfying soft constraints
- ★ Hard constraints are satisfied if feasible for  $x \in \mathcal{X}_{i|k}^{(r)}$  and  $x \in \mathcal{X}_{\mathbb{T}}^{(r)}$

## Affine uncertainty: Probabilistic tubes

Aside: the constraint

$$\frac{1}{N_s} \sum_{i=n+1}^{n+N_s} \Pr(f^T x_{k+i|k} + g^T u_{k+i|k} > h) < p$$

is satisfied for all  $n$  if  $\Pi$  and  $p_1, \dots, p_r$  satisfy

$$\frac{1}{N_s} \sum_{i=0}^{N_s-1} [p_1 \quad p_2 \quad \dots \quad p_r] \Pi^i e_1 < p$$

- ★ Fix  $\Pi$  and  $p_1, \dots, p_r$  offline and optimize  $\{\mathcal{X}_{i|k}^{(j)}\}$  online subject to constraints on:
  - (i). transition probabilities
  - (ii). probabilities of satisfying soft constraints
- ★ Hard constraints are satisfied if feasible for  $x \in \mathcal{X}_{i|k}^{(r)}$  and  $x \in \mathcal{X}_{\mathbb{T}}^{(r)}$

# Affine uncertainty: Tube constraints

For tractable implementation:

- 1 apply constraints to  $\{\mathcal{X}_{i|k}^{(j)}\}$  instead of  $\{\mathcal{S}_{i|k}^{(j)}\}$
- 2 invoke transition probability constraints via **inequality** constraints



Define  $\tilde{p}_{jm} = \sum_{l=1}^j p_{lm}$   $j, m = 1, \dots, r$ , and invoke constraints:

(a). on transition probabilities via

$$\Pr(x_{k+i+1|k} \in \mathcal{X}_{i+1|k}^{(j)} \mid x_{k+i|k} \in \mathcal{X}_{i|k}^{(m)}) \geq \tilde{p}_{jm} \quad i \leq N-1$$

$$\Pr(x_{k+i+1|k} \in \mathcal{X}_T^{(j)} \mid x_{k+i|k} \in \mathcal{X}_T^{(m)}) \geq \tilde{p}_{jm} \quad i \geq N$$

(b). on the probability of constraint satisfaction within  $\mathcal{X}_*^{(j)}$  via

$$\Pr(f^T x_{k+i+1|k} + g^T u_{k+i+1|k} > h \mid x_{k+i|k} \in \mathcal{X}_{i|k}^{(j)}) < p_j \quad i \leq N-1$$

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$$\Pr(x_{k+i+1|k} \in \mathcal{X}_{\mathbb{T}}^{(j)} \mid x_{k+i|k} \in \mathcal{X}_{\mathbb{T}}^{(m)}) \geq \tilde{p}_{jm} \quad i \geq N$$

(b). on the **probability of constraint satisfaction** within  $\mathcal{X}_{\star}^{(j)}$  via

$$\Pr(f^T x_{k+i+1|k} + g^T u_{k+i+1|k} > h \mid x_{k+i|k} \in \mathcal{X}_{i|k}^{(j)}) < p_j \quad i \leq N - 1$$

$$\Pr(f^T x_{k+i+1|k} + g^T u_{k+i+1|k} > h \mid x_{k+i|k} \in \mathcal{X}_{\mathbb{T}}^{(j)}) < p_j \quad i \geq N$$

# Affine uncertainty: Tube constraints

The constraint

$$\Pr(f^T x_{k+i|k} + g^T u_{k+i|k} > h) < p$$

is satisfied under (a) & (b) for all  $i$  if  $\Pi$  and  $p_1, \dots, p_r$  satisfy

$$[p_1 \ p_2 \ \cdots \ p_r] \Pi^i e_1 > p, \quad \forall i$$

and, additionally, if  $p_j$  and  $\tilde{p}_{jm}$  satisfy

$$\left. \begin{array}{l} p_j < p_{j+1} \\ \tilde{p}_{jm} \geq \tilde{p}_{j\ m+1} \end{array} \right\} \text{ for } j = 1, \dots, r-1$$

- ★  $p_j < p_{j+1}$   $\iff$  probability of satisfying soft constraint increases towards centre of tube
- ★  $\tilde{p}_{jm} \geq \tilde{p}_{j\ m+1}$   $\leftarrow$  always holds (since  $\mathcal{X}_*^{(1)} \subseteq \dots \subseteq \mathcal{X}_*^{(r)}$ )

# Affine uncertainty: Tube constraints

- Invoke the constraints on:
- (a). transition probabilities
  - (b). probabilities of satisfying soft constraints

using polytopic confidence regions for disturbance  $q$ , e.g.:

$$x_{k+i+1|k} \in \text{Co}\{\hat{A}^{(j)}x_k + \hat{B}^{(j)}u_k + \hat{d}^{(i)}, j = 1, \dots, l\} \text{ w.p. } \tilde{p}_{jm}$$

implies

$$\Pr(x_{k+i+1|k} \in \mathcal{X}_{i+1|k}^{(j)} \mid x_{k+i|k} \in \mathcal{X}_{i|k}^{(m)}) > \tilde{p}_{jm}$$

whenever

$$\text{Co}\{(\hat{A}^{(j)} + \hat{B}^{(j)}K_{\mathbb{T}})x_{i|k}^{(m,r)} + \hat{B}^{(j)}c_{k+i|k} + \hat{d}^{(j)}, j = 1, \dots, l\} \subseteq \mathcal{X}_{i+1|k}^{(m)}, r = 1, \dots, s$$

$$\text{where } \mathcal{X}_{i|k}^{(m)} = \text{Co}\{x_{i|k}^{(m,r)}, r = 1, \dots, s\}$$



finite set of linear constraints in the variables  $\mathbf{c}_k, \{x_{i|k}^{(m,j)}\}$

online optimization via QP

# Affine uncertainty: Approximate tubes

## Example

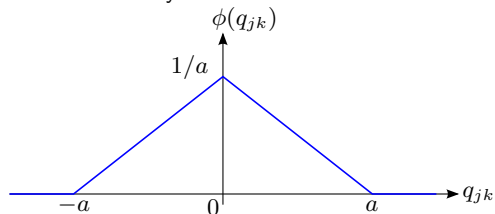
- ▷ Model parameters:

$$A^{(0)} = \begin{bmatrix} 1.2 & 0.1 \\ 0.1 & 1.26 \end{bmatrix} \quad A^{(1)} = \frac{10^{-2}}{a} \begin{bmatrix} -1 & -0.5 \\ -1 & 0.2 \end{bmatrix}, \quad A^{(2)} = \frac{10^{-2}}{a} \begin{bmatrix} -0.6 & 0.7 \\ -0.3 & 0.3 \end{bmatrix}$$

$$B^{(0)} = \begin{bmatrix} 0.5 \\ 0.21 \end{bmatrix}, \quad B^{(1)} = \frac{10^{-3}}{a} \begin{bmatrix} -1 \\ -2 \end{bmatrix}, \quad B^{(2)} = \frac{10^{-3}}{a} \begin{bmatrix} 2 \\ -9 \end{bmatrix}$$

$$d^{(1)} = \frac{1}{a} \begin{bmatrix} 0.1 \\ 0.01 \end{bmatrix}, \quad d^{(2)} = \frac{1}{a} \begin{bmatrix} 0.5 \\ 0.12 \end{bmatrix} \quad a = \sqrt{6}.$$

- ▷ Uncertainty distribution:



- ▷ Transition and constraint violation probabilities for  $r = 2$ :

$$\Pi = \begin{bmatrix} p_{11} & p_{12} \\ p_{21} & p_{22} \end{bmatrix} = \begin{bmatrix} 0.3 & 0.1 \\ 1 & 1 \end{bmatrix}$$

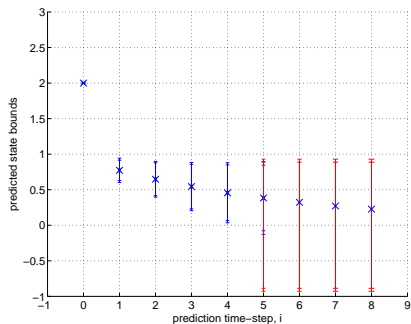
$$\begin{bmatrix} p_1 \\ p_2 \end{bmatrix} = \begin{bmatrix} 0.1 \\ 0.3 \end{bmatrix}, \quad p = 0.4$$



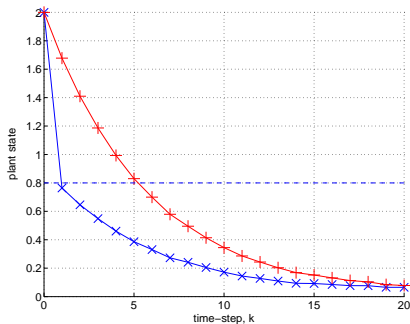
# Affine uncertainty: Approximate tubes

Constraint:  $\Pr\{x_1 > 0.8\} < 0.4$

Horizon:  $N = 5$



— :  $\{\text{Proj}_{x_1}(\mathcal{X}_i^{(j)}), i = 1, \dots, 5, j = 1, 2\}$   
 — :  $\{\text{Proj}_{x_1}(\mathcal{X}_T^{(j)}), j = 1, 2\}$

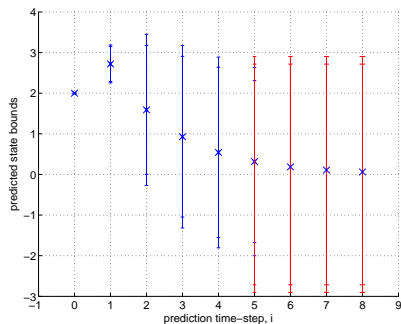


— :  $\mathbb{E}(x_1(k))$ , constrained MPC  
 — :  $\mathbb{E}(x_1(k))$ , unconstrained optimal

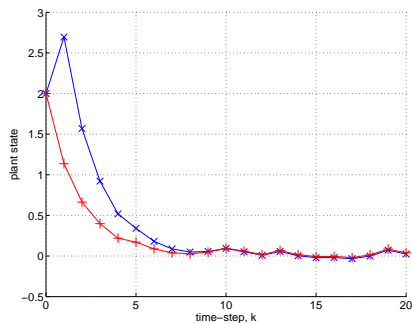
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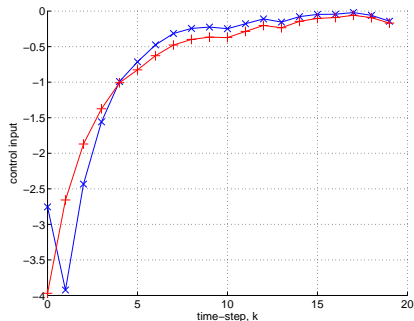
— :  $\{\text{Proj}_{x_2}(\mathcal{X}_i^{(j)}), i = 1, \dots, 5, j = 1, 2\}$   
 — :  $\{\text{Proj}_{x_2}(\mathcal{X}_T^{(j)}), j = 1, 2\}$



— :  $\mathbb{E}(x_2(k))$ , constrained MPC  
 — :  $\mathbb{E}(x_2(k))$ , unconstrained optimal

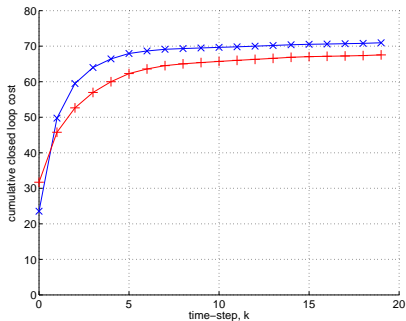
# Affine uncertainty: Approximate tubes

$\mathbb{E}(u(k))$ :



— : constrained MPC  
 — : unconstrained optimal

$\mathbb{E}(\|x(k)\|_Q^2 + \|u(k)\|_R^2)$ :



Target constraint violation rate: 0.4, actual: 0.2 (conservative!)

## Affine uncertainty: Exact tubes

## Affine uncertainty: Exact tubes

Invoke probabilistic constraints by sampling uncertainty distributions directly

[Tempo, Calafiore, Campi]

- ★ e.g. constraints:  $x(k) \in \mathcal{T}_0$  w.p. 1 &  $F(x(k), u(k), \mathfrak{D}) \leq 1$  w.p.  $p$  invoked via the constraints:

$$\forall x \in \mathcal{X}_k : \quad \frac{1}{N_\sigma} \sum_{j=1}^{N_\sigma} \sigma_j(x) \geq p, \quad \sigma_j(x) = \begin{cases} 1, & F(x, u(k), q_j) \leq 1 \\ 0, & F(x, u(k), q_j) > 1 \end{cases}$$

where  $\{x(0), \mathcal{X}_1, \mathcal{X}_2, \dots\}$  = robust uncertain tube given  $q \in \mathcal{Q} = \text{supp}(\mathfrak{D})$  w.p. 1

- ★ Bounds exist on  $N_\sigma$  to ensure satisfaction of  $\Pr\{F(x, u, \mathfrak{D}) \leq 1\} \geq p$  w.p.  $1 - \epsilon$

[Campi]

- ★ Avoids conservative parametric confidence bounds and computationally expensive numerical convolutions

# Affine uncertainty: Exact tubes

Assume:

(i)  $\mathcal{Q}$ : convex, bounded, polytopic

(ii)  $F(x, u, q)$ : convex in  $(x, u)$

[Prekopa, 1995]

then:  $\mathcal{X}_k = \text{Co}\{x^{(k,l)}, l = 1, \dots, r\}$ ,

and  $x(k) \in \mathcal{T}_0$  &  $\Pr\{F(x, u, \mathcal{D}) \leq 1\} \geq p$  invoked via

$$\forall l = 1, \dots, r : \quad \frac{1}{N_\sigma} \sum_{j=1}^{N_\sigma} \sigma_j(x^{(k,l)}) \geq p, \quad \sigma_j(x^{(k,l)}) = \begin{cases} 1, & F(x^{(k,l)}, u^{(k,l)}, q_j) \leq 1 \\ 0, & F(x^{(k,l)}, u^{(k,l)}, q_j) > 1 \end{cases}$$



Finite number of mixed integer linear constraints

[e.g. Blackmore, 2006]

# Application to terminal set calculation

Given r.p.i.  $\Omega$  set such that

$$F(x, K_{\mathbb{T}}x, \mathcal{D}) \leq 1 \text{ w.p. } p \quad \forall x \in \Omega$$

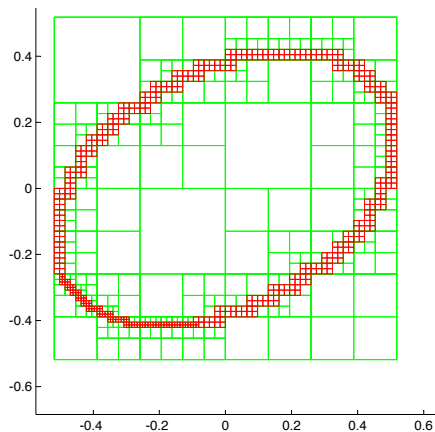
(i) cover  $\partial\Omega$  with boxes  $\Pi_i$

(ii) If:

$$\left. \begin{array}{l} f(x, K_{\mathbb{T}}x, \mathcal{D}) \in \Omega \text{ w.p. } 1 \\ F(x, K_{\mathbb{T}}x, \mathcal{D}) \leq 1 \text{ w.p. } p \end{array} \right\} \forall x \in \Pi_i$$

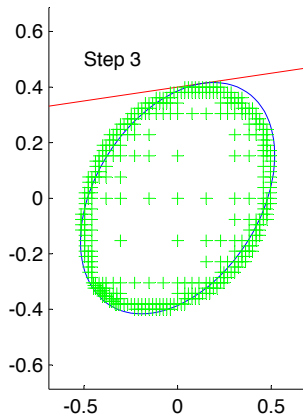
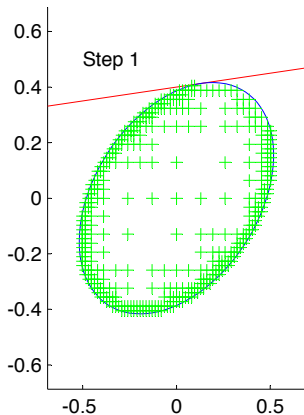
then  $\Pi_i \in \partial\hat{\Omega}$

(iii) set  $\Omega = \text{Co}\{\partial\hat{\Omega}\}$  and return to (i).



# Application to terminal set calculation

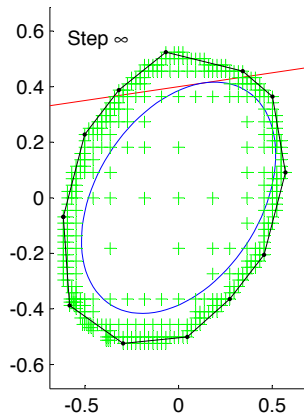
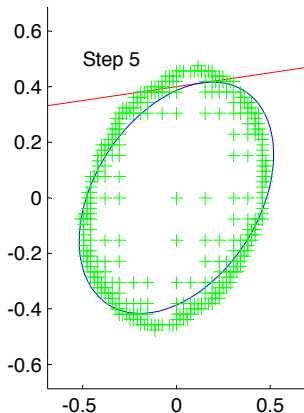
## Example





# Application to terminal set calculation

## Example



## Additive uncertainty: Exact tubes

# Additive uncertainty: problem formulation

Linear uncertain system:

$$\text{plant model} \quad x_{k+1} = Ax_k + Bu_k + w_k \quad x_k \in \mathbb{R}^n$$

$$\text{disturbance} \quad w_k = Dq_k, \quad q_k \in \mathcal{Q} \quad \mathcal{Q} \subset \mathbb{R}^m$$

$\{q_0, q_1, q_2, \dots\}$  assumed iid, with

- ▷  $q_{k,i}$ : zero-mean, independent, with known distributions:

$$\Pr\{q_{k,i} \leq \xi_i\} = \mathcal{F}_i(\xi_i), \quad i = 1, \dots, m$$

- ▷  $\mathcal{F}_i$ : right-continuous, with finitely many discontinuities and

$$\mathcal{F}_i(\xi) = \begin{cases} 1 & \text{for } \xi \geq \alpha_i \\ 0 & \text{for } \xi < -\alpha_i \end{cases}$$

↓

$$\mathcal{Q} \text{ compactly supported, } \mathcal{Q} = \{q : |q| \leq \alpha\}, \quad \alpha = \begin{bmatrix} \alpha_1 \\ \vdots \\ \alpha_m \end{bmatrix}$$

# Additive uncertainty: problem formulation

- Linear probabilistic constraints:

$$\Pr\{f_x^T x_k + g^T u_k \leq h\} \geq p$$

- Quasi-closed loop input predictions:  $u_{k+i|k} = K_{\mathbb{T}}x_{k+i|k} + c_{k+i|k}$ ,  
 $c_{k+i|k} = 0, \quad i \geq N$

State decomposition:  $x_{k+i|k} = z_{k+i|k} + e_{k+i|k} \quad \begin{cases} z_{k+i|k}: & \text{nominal} \\ e_{k+i|k}: & \text{uncertain} \end{cases}$

$$\begin{array}{ll} z_{k+i+1|k} = \Phi z_{k+i|k} + Bc_{k+i|k} & z_k|k = x_k \\ e_{k+i+1|k} = \Phi e_{k+i|k} + Dq_{k+i|k} & e_k|k = 0 \end{array}$$

$$\Phi = A + BK_{\mathbb{T}}$$

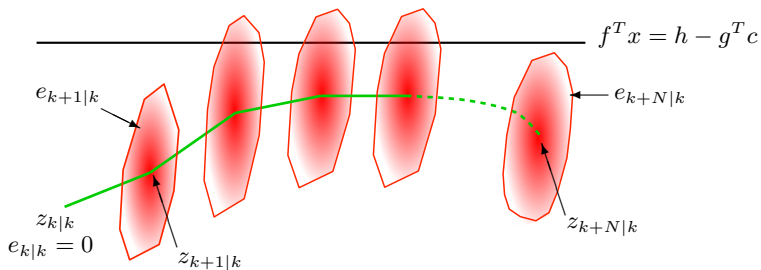
- Probabilistic constraints on predictions:

$$\Pr\{f^T x_{k+i|k} + g^T c_{k+i|k} \leq h\} \geq p$$

# Additive uncertainty: tubes

Future trajectories  $\{x_{k+i|k}, i = 0, 1, \dots\}$  belong to a tube

centred on  $\{z_{k+i|k}, i = 0, 1, \dots\}$ :



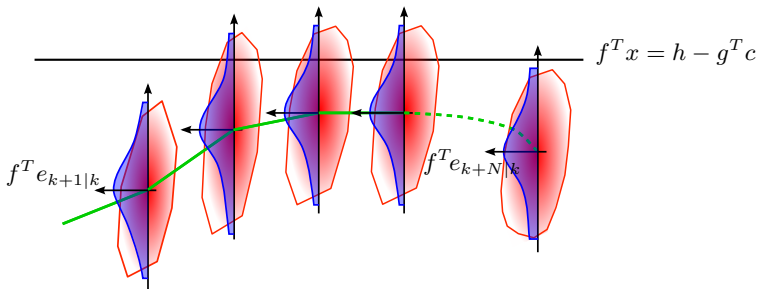
Compute the distributions of the projections  $f^T e_{k+i|k}$  directly

via a sequence of 1-dimensional convolutions

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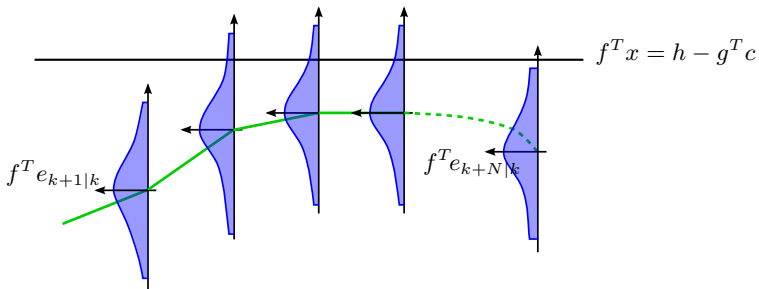
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Compute the distributions of the projections  $f^T e_{k+i|k}$  directly

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## Additive uncertainty: probabilistic constraints

Define  $\gamma_i$  as the minimum value such that

$$\Pr\{f^T e_{k+i|k} \leq \gamma_i\} = p$$

then  $\Pr\{f^T x_{k+i|k} + g^T c_{k+i|k} \leq h\} \geq p$  iff

$$f^T z_{k+i|k} + g^T c_{k+i|k} \leq h - \gamma_i$$

- Tightened linear constraints on nominal input/state predictions
- Given the distribution of  $\{w_0, w_1, w_2, \dots\}$ ,

compute  $\gamma_i$  for  $i = 1, 2, \dots$  using

$$\begin{aligned} f^T e_{k+i|k} &= f^T \Phi^{i-1} D q_{k|k} + \dots + f^T D q_{k+i-1|k} \\ &= \text{sum of independent, scalar r.v.'s} \end{aligned}$$

- $\gamma_i$  can be computed offline



# Additive uncertainty: numerical convolution

For scalar r.v.'s  $X, Y$  with densities  $f_X(x) : \Pr\{X \leq x\} = \int^x f_X dx$   
 $f_Y(y) : \Pr\{Y \leq y\} = \int^y f_Y dy$

the pdf of  $X + Y$  is the convolution

$$f_{X+Y} = f_X * f_Y$$

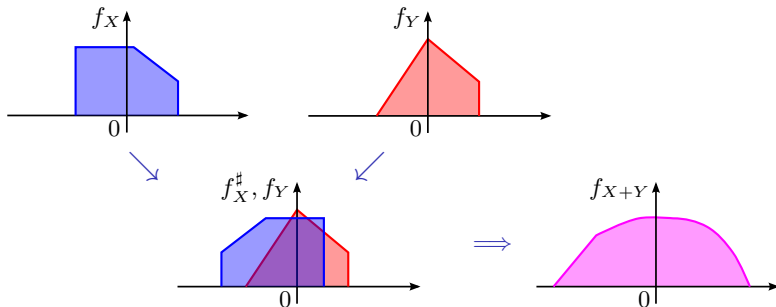
$f_{X+Y}$  can be approximated to required accuracy via discrete convolution

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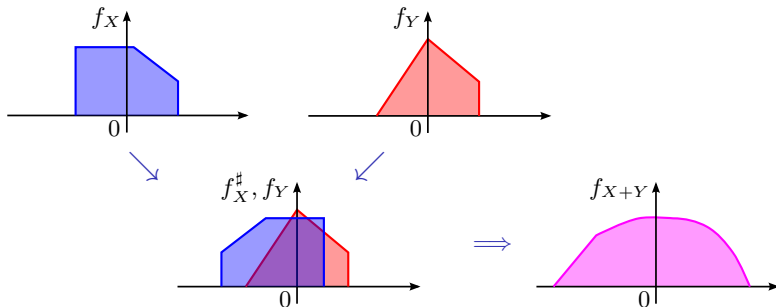
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the pdf of  $X + Y$  is the convolution

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$f_{X+Y}$  can be approximated to required accuracy via discrete convolution

# Additive uncertainty: numerical convolution

- Discretization of distribution functions  $F_{f^T \Phi^i Dq}(\cdot)$  on  $r$  intervals implies:
  - ★ approximation error:  $O(1/r^2)$  (e.g. trapezoidal integration)
  - ★ computation:  $O(r^2)$  multiplications/additions per convolution
- $\text{supp}(dF_{f^T \Phi^i Dq})$  increases monotonically with  $i$ , but is finite for all  $i$ .

N.B.  $\gamma_i$  is bounded  $\forall i$  (since  $\Phi$  is strictly stable)



e.g. Chebychev's one-sided inequality gives  $\gamma_i \leq \alpha \sqrt{f^T \Sigma_i f}$

where

$$\alpha^2 = p/(1-p)$$

$$\Sigma_1 = D\mathbb{E}(qq^T)D^T$$

$$\Sigma_{i+1} = \Phi \Sigma_i \Phi^T + D\mathbb{E}(qq^T)D^T, \quad i = 1, 2, \dots$$

# Additive uncertainty: recursive feasibility

Consider the  $i$ -step-ahead prediction at time  $k$ :

$$f^T e_{k+i|k} = f^T \Phi^{i-1} D q_k + f^T \Phi^{i-2} D q_{k+1} + \cdots + f^T D q_{k+i-1}$$



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at time  $k+1$ , this term has already been realized



Best bound on  $f^T e_{k+i|k+1}$  given information on  $q_k$  available at time  $k$ :

$$f^T e_{k+i|k+1} \leq a_{i-1} + f^T \Phi^{i-2} D q_{k+1} + \cdots + f^T D q_{k+i-1}$$

where

$$a_{i-1} = \max_{q \in \mathcal{Q}} f^T \Phi^{i-1} D q$$

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$\uparrow$   
 worst case bound

$\uparrow$   
 probabilistic bound



## Additive uncertainty: recursive feasibility

Consider the  $i$ -step-ahead prediction at time  $k$ :

$$f^T e_{k+i|k} = \underbrace{f^T \Phi^{i-1} D q_k}_{\text{worst case bound}} + \underbrace{f^T \Phi^{i-2} D q_{k+1} + \dots + f^T D q_{k+i-1}}_{\text{probabilistic bound}}$$

⇓

Best bound on  $f^T e_{k+i|k+1}$  with probability  $p$  given information available at  $k$ :

$$f^T e_{k+i|k+1} \leq \alpha_{i-1} + \gamma_{i-1} \quad \text{w.p. } p$$

where

$$\alpha_{i-1} = \max_{q \in \mathcal{Q}} f^T \Phi^{i-1} D q$$

and  $\gamma_{i-1}$  is the minimum value such that

$$\Pr\{f^T \Phi^{i-2} D q_{k+1} + \dots + f^T D q_{k+i-1} \leq \gamma_{i-1}\} = p$$



## Additive uncertainty: recursive feasibility

- Predictions at time  $k$  must ensure feasibility at  $k + 1, k + 2, \dots$
- Hence tighten constraints on nominal  $i$ -step-ahead prediction by  $\beta_i$ ,  
where  $\beta_i =$  maximum element of  $i$ th column of:

$$\begin{bmatrix} \gamma_1 & \gamma_2 & \gamma_3 & \gamma_4 & \dots \\ 0 & \gamma_1 + a_1 & \gamma_2 + a_2 & \gamma_3 + a_3 & \dots \\ 0 & 0 & \gamma_1 + a_1 + a_2 & \gamma_2 + a_2 + a_3 & \dots \\ 0 & 0 & 0 & \gamma_1 + a_1 + a_2 + a_3 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

Satisfaction of **probabilistic constraints** and **recursive feasibility** is ensured if

$$f^T z_{k+i|k} + g^T c_{k+i|k} \leq h - \beta_i$$

for  $i = 0, 1, 2, \dots$ , at each time  $k$

# Additive uncertainty: recursive feasibility

Properties of  $\beta_i$

$$(i) \beta_i = \gamma_1 + \sum_{j=1}^{i-1} a_j \quad \text{for } i = 1, 2, \dots$$



Largest element of each column lies on the diagonal:

$$\begin{bmatrix} \gamma_1 & \gamma_2 & \gamma_3 & \gamma_4 & \dots \\ 0 & \gamma_1 + a_1 & \gamma_2 + a_2 & \gamma_3 + a_3 & \dots \\ 0 & 0 & \gamma_1 + a_1 + a_2 & \gamma_2 + a_2 + a_3 & \dots \\ 0 & 0 & 0 & \gamma_1 + a_1 + a_2 + a_3 & \dots \\ \vdots & \vdots & \vdots & \vdots & \dots \end{bmatrix}$$

- ▷ Follows from  $\gamma_{i+1} \leq \gamma_i + a_i$
- ▷ Intuitively: future feasibility depends on worst case bounds on disturbances that have already been realized



$\{\beta_1, \beta_2, \beta_3, \dots\}$  is monotonically increasing

# Additive uncertainty: recursive feasibility

Properties of  $\beta_i$

$$(i) \beta_i = \gamma_1 + \sum_{j=1}^{i-1} a_j \quad \text{for } i = 1, 2, \dots$$



Largest element of each column lies on the diagonal:

$$\begin{bmatrix} \gamma_1 & \gamma_2 & \gamma_3 & \gamma_4 & \dots \\ 0 & \gamma_1 + a_1 & \gamma_2 + a_2 & \gamma_3 + a_3 & \dots \\ 0 & 0 & \gamma_1 + a_1 + a_2 & \gamma_2 + a_2 + a_3 & \dots \\ 0 & 0 & 0 & \gamma_1 + a_1 + a_2 + a_3 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

- ▷ Follows from  $\gamma_{i+1} \leq \gamma_i + a_i$
- ▷ Intuitively: future feasibility depends on worst case bounds on disturbances that have already been realized



$\{\beta_1, \beta_2, \beta_3, \dots\}$  is monotonically increasing

# Additive uncertainty: recursive feasibility

Properties of  $\beta_i$

(ii)  $\lim_{i \rightarrow \infty} \beta_i \leq \bar{\beta}_\nu$ , where  $\bar{\beta}_\nu$  is defined for any integer  $\nu \geq 1$  by

$$\bar{\beta}_\nu = \gamma_1 + \sum_{j=1}^{\nu-1} a_j + \frac{\rho^\nu}{1-\rho} \|g\|_S$$

where  $\|g\|_S = \sqrt{g^T S g}$  and  $\rho, S$  satisfy

$$\begin{aligned} \max_{q \in \mathcal{Q}} \|Dq\|_{S^{-1}} &\leq 1 \\ \Phi S \Phi^T &\leq \rho^2 S, \quad \rho \in (0, 1) \end{aligned}$$

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Follows from

$$\star a_j = \max_{q \in \mathcal{Q}} g^T \Phi^j Dq \leq \max_{\|v\|_{S^{-1}} \leq 1} g^T \Phi^j v \leq \|\Phi^j\|_S \|g\|_S$$

$$\star \max_{q \in \mathcal{Q}} g^T \Phi^j Dq \leq \rho \|\Phi^{j-1}\|_S \|g\|_S$$

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$$\bar{\beta}_\nu - \frac{\rho^\nu}{1-\rho} \|g\|_S \leq \lim_{i \rightarrow \infty} \beta_i \leq \bar{\beta}_\nu$$

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$\lim_{i \rightarrow \infty} \beta_i$  may be determined to any desired accuracy  $\epsilon > 0$

if  $\nu$  is chosen to be sufficiently large that  $\frac{\rho^\nu}{1-\rho} \|g\|_S < \epsilon$

## Additive uncertainty: SMPC algorithm

Constraints on predicted  $\{z_{k+i|k}, c_{k+i|k}\}$  in MPC optimization at time  $k$ :

$$f^T z_{k+i|k} + g^T c_{k+i|k} \leq h - \beta_i, \quad i = 0, \dots, N-1$$

$$z_{k+N|k} \in \mathcal{S}_\nu$$

where the **terminal constraint** set  $\mathcal{S}_\nu$  is defined by

$$\mathcal{S}_\nu = \{z : f^T \Phi^{i-N} z \leq h - \beta_i, \quad i = N, \dots, \nu - 1$$

$$f^T \Phi^{i-N} z \leq h - \bar{\beta}_\nu, \quad i = \nu, \nu + 1, \dots\}$$



These constraints are sufficient to ensure recursive feasibility since

$$\mathcal{S}_\nu \supseteq \mathcal{S}_\infty$$

where  $\mathcal{S}_\infty$  is the **maximal admissible set**:

$$\mathcal{S}_\infty = \{z : f^T \Phi^{i-N} z \leq h - \beta_i, \quad i = N, N+1, \dots\}$$



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$$z_{k+N|k} \in \mathcal{S}_\nu$$

Properties of  $\mathcal{S}_\nu$ :

(i)  $\mathcal{S}_\nu$  is compact and non-empty iff  $h \geq \bar{\beta}_\nu$  (assuming  $(\Phi, g)$  observable)

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$$z_{k+N|k} \in \mathcal{S}_\nu$$

Properties of  $\mathcal{S}_\nu$ :

(ii)  $\mathcal{S}_\nu$  is finitely determined:

$$\mathcal{S}_\nu = \{z : f^T \Phi^{i-N} z \leq h - \beta_i, \quad i = N, \dots, \nu - 1$$

$$f^T \Phi^{i-N} z \leq h - \bar{\beta}_\nu, \quad i = \nu, \nu + 1, \dots, \nu + N^*\}$$

- ★  $N^*$  can be computed using e.g. [Gilbert & Tan, 1991]
- ★  $\mathcal{S}_\nu$  is assumed to be compact

# Additive uncertainty: SMPC algorithm

Algorithm:

*Offline* Determine parameters  $\gamma_1, \beta_1, \dots, \beta_{\nu-1}, \bar{\beta}_\nu, N^*$   
defining the recursively feasible probabilistic constraints

*Online* At each time  $k = 0, 1, \dots$ :

1. Obtain the current state  $x_k$
2. Solve the QP:

$$\begin{aligned} \mathbf{c}_k^* &= \arg \min_{\mathbf{c}} V(x_k, \mathbf{c}) \\ &\text{subject to } f^T z_{k+i|k} + g^T c_{k+i|k} \leq h - \beta_i, \quad i = 0, \dots, N-1 \\ &\quad z_{k+N|k} \in \mathcal{S}_\nu \end{aligned}$$

3. Set  $u_k = K_{\mathbb{T}} x_{k+i|k} + c_{k+i|k}^*$ .

# Additive uncertainty: State estimation example

- Output feedback MPC formulation:

$$\text{plant model} \quad x_{k+1} = Ax_k + Bu_k + Dw_k$$

$$\text{output} \quad y_k = Cx_k + Fv_k$$

disturbance  $w_k$  and measurement noise  $v_k$  iid & compactly supported

- Linear observer:

$$\text{estimate update} \quad \hat{x}_{k+1} = A\hat{x}_k + Bu_k + L(y_k - \hat{y}_k)$$

$$\text{with} \quad \hat{y}_k = C\hat{x}_k$$

state estimate:  $\hat{x}$ , observer gain:  $L$  such that  $(A - LC)$  is strictly stable

- Estimation error  $\epsilon_k := x_k - \hat{x}_k$  has dynamics:

$$\epsilon_{k+1} = (A - LC)\epsilon_k + Dw_k + Fv_k$$

$\epsilon_0$ : r.v. with known & compactly supported distribution and  $\epsilon_0 \in \Pi_0$

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# Additive uncertainty: State estimation example

## Model parameters

$$A = \begin{bmatrix} 1.6 & 1.1 \\ -0.7 & 1.2 \end{bmatrix}, \quad B = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \quad C = [0.9 \quad 0.2], \quad D = I, \quad F = 1$$

- Noise & uncertainty derived from truncated normal distributions:

$$v_k \sim \{ \mathcal{N}(0, 1/24^2) \quad \text{truncated so that} \quad |v_k| \leq 0.12 \}$$

$$w_k \sim \{ \mathcal{N}(0, I/24^2) \quad \text{truncated so that} \quad \|w_k\|_\infty \leq 0.12 \}$$

$$\epsilon_0 \sim \{ \mathcal{N}(0, I/24^2) \quad \text{truncated so that} \quad \|\epsilon_0\|_\infty \leq 0.12 \}$$

- Probabilistic state constraints:

$$\eta_x^T = [\pm 1 \quad \pm 0.3], \quad \eta_u = 0, \quad h = 1.5, \quad p = 0.8.$$

- $$\left. \begin{array}{l} K_{\mathbb{T}} = -[1.03 \quad 1.07]: \quad \text{unconstrained LQ-optimal} \\ L = -[0.83 \quad 1.22]: \quad \text{steady Kalman filter gain} \end{array} \right\} \implies |\lambda_{\max}(\Psi)| = 0.12$$

# Additive uncertainty: State estimation example

## Prediction parameters

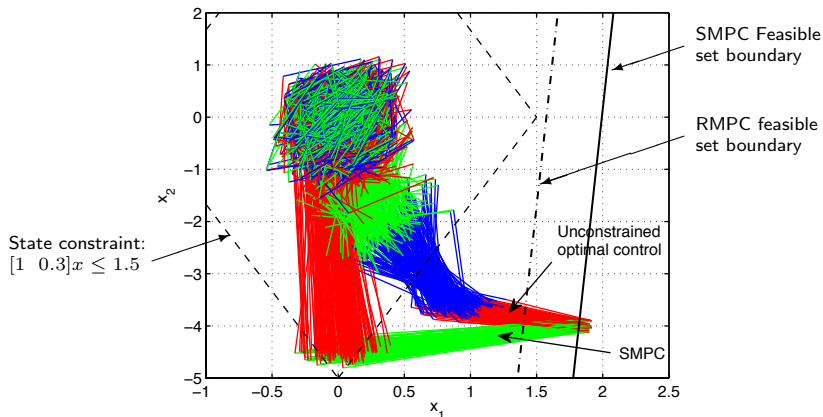
- Dof in predictions:  $N = 6$
- Horizon for transients:  $\nu = 13$
- Probabilistic bounds  $\gamma_{0|k}$  computed for  $k = 0, \dots, \nu - 1$   
using rectangular integration with grid spacing  $10^{-4}$
- Terminal sets: non-empty for  $h \geq \bar{\beta} = 1.31, 1.20$   
finitely determined, with  $N_k^* = 10$  for all  $k$

## Simulation parameters

- $10^4$  realizations of disturbance and noise sequences
- $x_0 = \hat{x}_0 + \epsilon_0$ , for  $10^4$  realizations of  $\epsilon_0$   
and fixed  $\hat{x}_0 = (1.8, -4.0)$



# Additive uncertainty: Example

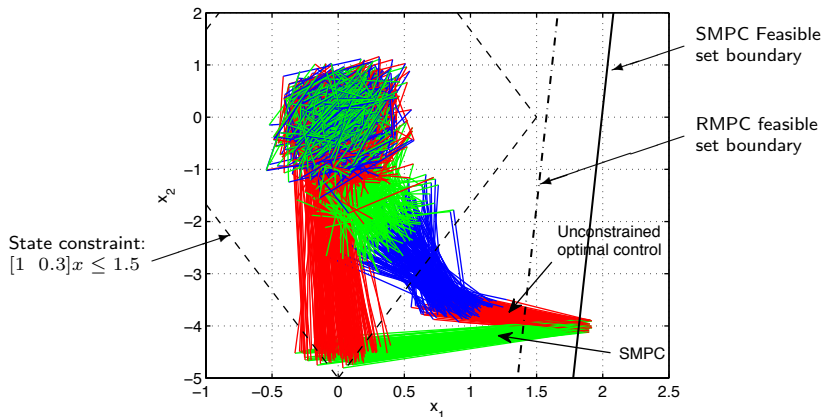


$\hat{x}_0 = (1.8, -4.0)$ :

Constraint violation frequency:

	target	observed (at $k = 1$ )
SMPC	20%	20.0%
LQ-optimal	—	100%

# Additive uncertainty: Example

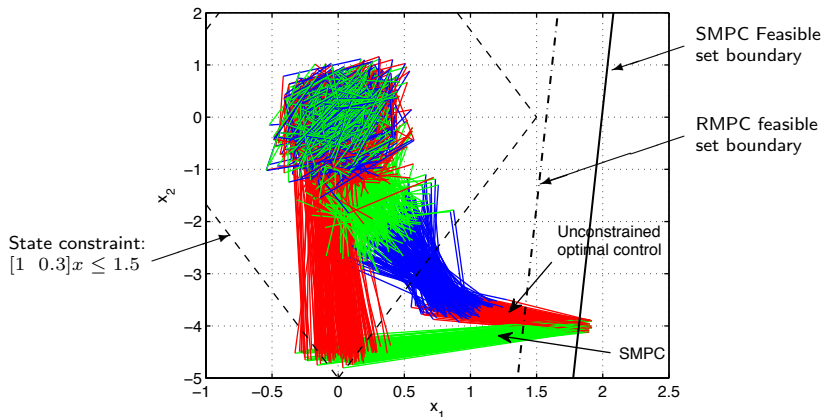


$$\hat{x}_0 = (1.8, -4.0):$$

Cumulative cost	$\sum_{k=0}^{12} \mathbb{E}(\ x_k\ _Q^2 + \ u_k\ _R^2):$	SMPC	3.56
		LQ-optimal	2.38

(RMPC infeasible)

# Additive uncertainty: Example



$$\hat{x}_0 = (1.35, -4.0):$$

Cumulative cost	$\sum_{k=0}^{12} \mathbb{E}(\ x_k\ _Q^2 + \ u_k\ _R^2):$	SMPC	1.72
		RMPC	2.42

# Conclusions

- ① Example: fatigue control
- ② Stochastic MPC: basic formulations
  - ▶ Probabilistic constraints & recursive feasibility
  - ▶ Performance costs and stability analyses
- ③ Implementation
  - ▶ Affine model uncertainty: approximate and exact tubes
  - ▶ Additive model uncertainty: exact tubes